Who owns the interface?

Linux Plumbers Conference

Portland, OR, USA; 18 Sep 2008

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http://www.kernel.org/pub/linux/docs/manpages

git://git.kernel.org/pub/scm/docs/man-pages/man-pages.git

http://linux-man-pages.blogspot.com

Overview

- What does it mean to own the interface?
- Who could claim to own the interface?
- Arguments for (*pros*) and against (*cons*) each claim
- Problems resulting from the Linux ownership model
- Conclusion/Discussion

What does it mean to “own the interface”?

- Who determines what gets included in the interface, and what doesn’t?
- How is “in the interface” even defined?
  - By the implementation (code below the interface)?
  - By the use cases (code above the interface)?
  - By a specification?

Traditional Interface Ownership Model

- (Strong) degree of centralization
- Interface owned by vendor
  - e.g., AT&T's SVID
- Or by a standards body
  - e.g., POSIX.1

But, who does own the Linux interface?

Candidates

- Austin
- Kernel developers
- Users
- glibc
- LSB
- LTP
- man-pages

Is it the Austin Group?

Pros

- They define standard (POSIX.1/SUS) that describes much of the Unix API
- Linux kernel (and glibc) strive to conform to standard

Cons

- Surprisingly long list!

The Austin Group - Cons

- Linux doesn't implement all POSIX.1 interfaces
- Many Unix/Linux interfaces not specified in POSIX.1
  - *e.g., setresuid(), mincore(), brk(), flock(), settimeofday()*
- Linux provides extensions
  - e.g., *inotify*, *epoll*, `/proc`, capabilities, extended attributes, NUMA, scheduler affinity, various process resource limits (*setrlimit()*), *sendfile()*, etc.

The Austin Group - Cons (2)

- Linux may deliberately violate the standard
  - e.g., *link()* to a symlink
- Sometimes, an implementation accidentally violates the standard
  - e.g., successful *sched_setscheduler()* should return previous scheduling policy, not 0
  - Define non-conformant details as being “part of the Linux interface”?

The Austin Group - Cons (3)

- SUS/POSIX.1 leaves many things unspecified
  - Examples:
    - Does *select()* modify its *timeout* argument?
    - Does *signal()* provide reliable semantics?
    - Do *setitimer()* and *alarm()* interact?
    - What is maximum size of a Unix domain datagram?
    - What are precise semantics of *vfork()*?
    - What is resolution of system clock?
    - Does *longjmp()* restore signal mask to state at time of *setjmp()*?
- Left to implementation to define (or not care about) them

The Austin Group?

- Many parts of Linux interface clearly **not** owned by Austin

Is it the Kernel Developers?

Pros

- *Surely* it is the kernel developers?
- After all, they implement the interface!

The Kernel Developers - Cons

- What about difference between implementation and intention?
  - What happens if the kernel-user interface has bugs?
    - Kernel developers write a *lot* of bugs in interfaces
  - Unforeseen uses of interface (see “Users” later)
- glibc mediates between kernel and user
  - e.g., syscall wrapper may provide different behavior from raw syscall
  - More on glibc later

The Kernel Developers – Cons (2)

- If implementation defines interface, how can user know what definition is?
- Read the source???
- Wrong answer for many reasons...
  - Takes too long
  - Doesn't tell us whether a feature is intended or a bug
  - More on reading source later...

Is it the Users?

Cons:

- How could it **possibly** be the users? They didn't write the interface!

Can there be any pros?

The Users - Pros

Bugs!

- Suppose an interface contains a bug
- What is “correct” definition of interface?
- The intention? or the implementation?
- Should we fix the bug?
- What if users already program around the bug?
  - Maybe better not fix…
  - Unless we can tell users they suck…
    - (i.e., intended behavior was documented)
  - Then, arguably, users have defined part of interface

The Users – Pros (2)

Reading the source, revisited

- Creates a “tight” specification
- Suppose implementer has some detail of interface behavior (“x”) that shouldn't be fixed in stone
  - maybe want to change it in future
- Don't want users to rely on detail “x” in their code
- Making users read the source:
  - Reveals all details of interface
  - Provides no warning against relying on detail “x”
  - Some users will write code that uses “x”
  - Again, users have defined part of interface

Is it the glibc Developers?

- Raw kernel syscall interface is like a newly built house: you want someone make it livable before you move in
- glibc mediates between kernel and user, providing syscall wrappers

Glibc - Pros

- Some wrappers do significant work on top of system calls
  - *stat(), readv(), writev(), pselect(), mq_getattr() / mq_setattr()*
- glibc implements many functions not based on syscalls
  - To users, these are part of Linux interface

Glibc - Cons

- See kernel “cons” -- bugs, unintended uses of API, etc
- For many syscalls, wrapper is trivial
  - glibc is transparently exposing kernel interface
- glibc doesn't provide wrappers for every system call
- What gets wrapped or not is a little arbitrary:
  - Did glibc folk notice new kernel interface?
  - Did they think it was worth wrapping?
- Sometimes, glibc developers deliberately choose not to wrap a syscall
  - *gettid()* (though thread IDs are needed by some syscalls)

Glibc – Cons (2)

- Many kernel interfaces other than syscalls
- Those interface are *not* mediated by glibc, e.g.:
  - `/proc`
  - sysfs
  - *ioctl()*
  - netlink

Is it LSB?

Pros

- LSB defines an ABI standard for Linux

Cons

- Many interfaces not specified (sometimes deliberately)
- Parts of some interfaces are deliberately unspecified
- LSB is largely just standardizing the implementation provided by kernel and glibc developers

Linux Test Project (LTP)?

Pros

- Tests embody a specification, in code

Cons

- Test coverage is not complete, and excludes glibc
- Tests may themselves be buggy
- LTP tests usually only added (well) after syscall is added to kernel

Is it me?

- The *man-pages* project documents kernel and glibc interfaces
- Documentation is contract between kernel (/ glibc) and user(?)

man-pages - Pros

- Documentation can describe the developer's *intention*
- Provides reference for determining where implementation deviates from intention
  - i.e., is this a bug?
- Documentation can loosen the specification:
  - Can say things like: you *can* do “x”, but if you do, the results are undefined
  - (cf. “Tight” specification that results from reading source)

man-pages - Cons

- Many things remain undocumented.
  - Does that mean that they are not part of the interface?
- Sometimes implementation is right and documentation is wrong :-(

Summary of Ownership

- Many groups have claims to ownership
- Some validity in each claim
- No group can claim exclusive ownership
- Distributed ownership is source of some problems

Problems resulting from distributed ownership

Many of the problems arise from a single point...

- **How do we even know when an interface has been added or changed?**
  - We == kernel developers, glibc developers, userland programmers, testers, LSB, *man-pages*

Problems resulting from distributed ownership

Consequences for documentation

- Documentation may be late (i.e., after implementation)
- Poor documentation
  - (esp. if implementer was not involved)
- No documentation

Problems resulting from distributed ownership

Consequences for testing

- Late and insufficient testing
- Insufficient pre-release testing →
- many bugs in released interfaces
  - *epoll, timerfd, utimensat(), signalfd()*

Problems resulting from distributed ownership

Consequences for interface design

- Insufficient design review before release
- Inconsistent interfaces
  - Rounding of args for *mlock()* and *remap_file_pages()*
- Poorly designed interfaces
  - Dnotify
- Design mistakes
  - *epoll_create()* nowadays ignores *size* arg

Problems resulting from distributed ownership

- Design mistakes
- Capabilities
- Divide *root* into many distinct pieces

  CAP_AUDIT_CONTROL CAP_AUDIT_WRITE CAP_CHOWN CAP_DAC_OVERRIDE CAP_DAC_READ_SEARCH CAP_FOWNER CAP_FSETID CAP_IPC_LOCK CAP_IPC_OWNER CAP_KILL CAP_LEASE CAP_LINUX_IMMUTABLE CAP_MAC_ADMIN CAP_MAC_OVERRIDE CAP_MKNOD CAP_NET_ADMIN CAP_NET_BIND_SERVICE CAP_NET_BROADCAST CAP_NET_RAW CAP_SETFCAP CAP_SETGID CAP_SETPCAP CAP_SETUID **CAP_SYS_ADMIN** CAP_SYS_BOOT CAP_SYS_CHROOT CAP_SYS_MODULE CAP_SYS_NICE CAP_SYS_PACCT CAP_SYS_PTRACE CAP_SYS_RAWIO CAP_SYS_RESOURCE CAP_SYS_TIME CAP_SYS_TTY_CONFIG
- Great! But which one do I (an implementer) use?
- Ahh! I know!
- `CAP_SYS_ADMIN`, the new *root*, 180 uses in 2.6.27-rc

Concluding thoughts

- Interfaces are contracts
- Cast in stone
- We live with them “forever”
- So: need to get them right, at the beginning
- Getting things right:
  - Requires some degree of planning and coordination
  - Probably more than we currently do.
- Linux may be evolution, but intelligent design might sometimes get us there better and faster

Discussion / Questions

- How do we know when an interface has been changed or added?

http://userweb.kernel.org/~mtk/papers/lpc2008/who_owns_the_interface.pdf